A Signal Temporal Logic Motion Planner for Bird Diverter Installation Tasks With Multi-Robot Aerial Systems

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ABSTRACT This paper addresses the problem of task assignment and trajectory generation for installing bird diverters using a fleet of multi-rotors. The proposed solution extends our previous motion planner to compute feasible and constrained trajectories, considering payload capacity limitations and recharging constraints. Signal Temporal Logic (STL) specifications are employed to encode the mission objectives and temporal requirements. Additionally, an event-based replanning strategy is introduced to handle unforeseen failures. An energy minimization term is also employed to implicitly save multi-rotor flight time during installation operations. The effectiveness and validity of the approach are demonstrated through simulations in MATLAB and Gazebo, as well as field experiments carried out in a mock-up scenario.

INDEX TERMS Aerial systems: applications, formal methods in robotics and automation, multi-robot systems, task and motion planning.

I. INTRODUCTION
Power lines play a crucial role in providing energy to millions of people and are considered vital civil infrastructure in any country. To enhance network reliability and minimize power outages, electricity supply companies invest significant resources in inspection and maintenance operations [1]. Among these activities, the installation of bird diverters on power lines (see Figure 1) is essential to mitigate the risk of bird collisions [2] and improve their visibility [3]. Bird mortality caused by power line collisions is a significant concern, particularly in areas with diverse bird populations or during migratory seasons. It has been estimated that power line collisions affect approximately 350 bird species [4], resulting in millions of bird deaths annually worldwide [5], [6].

To address this issue, various types of bird diverters have been developed, including active and passive designs. Active bird diverters utilize wind-driven components, while passive diverters, such as helical objects made of plastic or aluminum, are attached to power cables to serve as visual markers (see Figure 1). Additionally, alternative techniques, such as visual and auditory deterrents, have been developed to mitigate bird collisions. Visual deterrents employ markers or reflective materials to enhance visibility and assist birds in avoiding power lines, while auditory deterrents emit sounds that discourage birds from approaching. However, the effectiveness of these techniques can vary depending on bird species and local conditions. Integrating these techniques with other mitigation strategies provides a multi-faceted approach to enhance bird safety and maintain a reliable power supply [7]. However, the current method of using manned helicopters and experienced crews for the installation of bird diverters has drawbacks. Firstly, it is time-consuming, as power lines are often located in difficult-to-access areas. Secondly, it poses safety risks due to installations being performed at heights on active lines.

Unmanned Aerial Vehicles (UAVs) offer a promising solution to automate bird diverter installation [8], [9], [10]. They
can operate continuously over long distances and be equipped with lightweight manipulation devices for autonomous installation operations [8], [11], [12], [13], [14]. However, the limited battery and payload capacity of individual UAVs necessitate the use of multiple UAVs to expedite the process, respond to unforeseen events and cover large-scale scenarios.

Planning for a multi-UAV team presents challenges such as scheduling battery recharging and diverter installation, ensuring collision-free trajectories, considering vehicle dynamics and energy consumption models, among other concerns.

Therefore, advanced task and motion planning techniques are required to enable bird diverter installation using multi-UAV teams while meeting safety requirements and mission objectives. Temporal Logic (TL) can serve this purpose by providing a mathematical framework that combines natural language commands with temporal and Boolean operators [15], [16]. In particular, Signal Temporal Logic (STL) [17] is equipped with a metric called robustness, which not only evaluates whether the system execution meets requirements but also quantifies the extent to which these requirements are fulfilled. This leads to an optimization problem that aims to maximize the robustness score, thereby providing the best feasible trajectory while satisfying the desired specifications.

This paper proposes a motion planner for multi-UAV systems that leverages STL specifications for the installation of bird diverters on power lines. The mission requirements are encoded as an STL formula, and a nonlinear non-convex max-min optimization problem is formulated to maximize the robustness score. To handle the complexity of this nonlinear optimization, a hierarchical approach is adopted. First, a Mixed-Integer Linear Programming (MILP) problem is solved, and the resulting solution is then fed into the final STL optimizer.

A. RELATED WORK

This paper addresses a complex multi-UAV motion problem that requires considering vehicle dynamics, collision avoidance, limited mission time, and payload capacity. Sampling-based planners [18] like the well-known Rapidly-exploring Random Tree (RRT) and its variants can compute collision-free trajectories in high-dimensional configuration spaces within reasonable time. Extensions have been proposed to minimize energy consumption and handle vehicle dynamics constraints [19], [20], although their convergence is not proven. However, RRT-based methods do not inherently handle payload capacity restrictions and struggle with complex nonlinear problems or multi-robot settings. To achieve real-time performance, motion primitives have been explored for optimal trajectory planning of UAVs [21], [22]. Additionally, the problem often transforms into a combinatorial optimization challenge, where an exhaustive search becomes computationally infeasible. In such cases, heuristic methods are commonly employed to simplify the complexity and identify the most viable solution [23], [24], [25], [26].

When the mission requires visiting multiple locations, optimization approaches based on the Vehicle Routing Problem (VRP) formulation have been proposed [27], [28]. Since VRP variants are NP-hard combinatorial problems, many works propose heuristic approaches sacrificing optimality [29], [30]. While these methods can handle UAV capacity and mission time constraints, they face challenges in obstacle avoidance and multi-vehicle dynamics [31]. Hierarchical approaches utilizing Probabilistic Roadmaps (PRMs) and kinodynamic Stable Sparse RRT (SST) methods have been investigated to generate collision-free and minimum-time flight trajectories [32], [33]. However, the problems discussed do not consider UAV payload capacity or multi-UAV settings.

Alternatively, with the advancement of computational resources onboard UAVs and recent progress in efficient numerical methods, receding horizon approaches for optimal control have emerged as a viable solution for addressing multi-UAV trajectory planning [34]. Centralized methods for non-convex optimization [35] and distributed Model Predictive Control (MPC)-based algorithms [36], [37] have been proposed. Others have employed knowledge transfer techniques between agents to achieve collision avoidance and safety [38]. While these methods efficiently compute safe and optimal trajectories for multiple UAVs, they do not incorporate high-level mission specifications (e.g., time interval or ordering constraints) or task planning into the problem formulation.

Moreover, literature includes publications that enhance motion planning for multi-robot systems using formal specification languages [17], [39]. The authors in [40] propose extracting local specifications assigned to specific robots to address multi-robot trajectory planning. Motion capabilities of each robot are represented using a transition system, which may become computationally expensive as the number of agents and tasks increase. On the other hand, Control Barrier Functions (CBFs) [41] have shown promise in efficient motion planning, but they are limited to a fragment of the STL formula, restricting their application to simple scenarios [42]. In [43], the authors present robust algorithms for multi-robot coordination that encode high-level temporal logic specifications into an MILP problem. However, their sequential multi-robot RRT algorithm leads to...
suboptimal trajectories. A reinforcement learning approach for multi-agent systems has been proposed in [44], focusing on finite abstractions and deterministic systems. Nevertheless, methodological advances are still needed as it poses computational challenges for complex STL formulae. In [45], the authors propose a mathematical programming-based method for MPC using STL specifications. Nonetheless, its scalability is limited due to exponential complexity in the number of binary variables. In [46], a scalable algorithm for multi-agent optimal control with TL is presented, but it lacks the ability to handle concrete timing requirements and agent counting.

### B. CONTRIBUTIONS

This paper introduces an innovative approach to task and motion planning for the installation of bird diverters on power lines using a team of multi-rotor UAVs. The proposed method leverages STL to generate optimal trajectories that are dynamically feasible, considering vehicle dynamics and velocity and acceleration limits. These trajectories fulfill various mission specifications, including collision avoidance between UAVs and the environment, adherence to UAV payload capacity limits, and reaching all installation targets within the given mission time while ensuring sufficient time for diverter installation and refilling operations. The problem results in a complex nonlinear non-convex max-min optimization problem, which is addressed through a novel hierarchical approach.

Building upon our previous research [47], the work is extended to specifically address the installation of bird diverters, incorporating payload capacity limitations and considering the presence of refilling stations for long-endurance operations. However, finding an optimal solution in a reasonable time frame can be challenging due to the solvers’ tendency to get stuck in local optima based on the initial guess [48], [49]. Therefore, the work [47] is also extend by introducing a novel procedure to compute the initial guess solution for the optimization problem. Additionally, our approach incorporates an event-based replanning strategy to handle unforeseen failures. Whenever a UAV fails, a new plan is computed online using the motion planner for a backup drone. This ensures uninterrupted mission progress. Lastly, to implicitly save multi-rotor flight time during installation operations, the proposed method also integrates an energy minimization term. The core idea is to determine an optimal cruising speed for each UAV, minimizing energy consumption, and generating trajectories that closely follow these nominal speeds. The main contributions of this paper are summarized as follows:

- The multi-UAV planning problem for installing diverters on power lines is formulated in Section II. We extract STL specifications (see Section III) and employ them to establish an optimization problem that integrates task and motion planning. This results in dynamically feasible trajectories that satisfy safety requirements, including obstacle avoidance and mutual safety distances, while ensuring mission fulfillment.
- The proposed STL optimization problem (Section IV) provides global optimal solutions. However, its nonlinear non-convex max-min nature presents challenges for solution methods [48], [49]. To address this, we propose a hierarchical approach where an MILP approach (Section IV-B) generates the initial guess. This MILP approach works with simplified constraints, neglecting obstacle avoidance, safety requirements, vehicle dynamics, and time mission specifications. It outputs the sequence of target regions and corresponding refilling operations for each UAV. Motion primitives [47] are then utilized to approximate UAV dynamics and generate feasible trajectories (position, velocity, and acceleration) between each pair of points. These trajectories serve as the initial guess for the global STL planner, aiding in its convergence towards solutions that meet all requirements.
- An event-based replanning strategy (Section IV-C) is introduced to handle unforeseen UAV failures, making the method applicable to scenarios with dynamic conditions. Additionally, an extension to the proposed planner (Section IV) is presented in the form of an energy-aware planner (Section IV-D). This enhanced planner includes an energy minimization term to implicitly extend the endurance of the multi-rotors during the mission.
- Numerical simulations in MATLAB (Sections V) assess the overall performance of the method in terms of mission specification fulfillment and the effectiveness of the novel initialization procedure for the STL optimization. Furthermore, Gazebo simulations and field experiments (Section V-D) conducted in a mock-up setting demonstrate the validity and feasibility of the method for real scenarios. The Discussion section (Section VI) provides a critical analysis of the paper’s findings, exploring their implications, limitations, and potential future extensions. Conclusions are discussed in Section VII.

### II. PROBLEM DESCRIPTION

This paper focuses on the problem of installing bird diverters on power lines using a team of $\delta$ UAVs. The work presented here forms a part of the AERIAL-CORE European project\(^1\). Figure 2 illustrates the scenario setup. The objective is to visit specific target regions along the upper cables between consecutive towers for installation purposes [9]. These regions serve as simplified representations of 3D spaces where diverters can potentially be installed. Once a UAV reaches a region, it is assumed that an onboard low-level controller handles the installation operation, e.g., [8], [10], [12]. The UAVs are assumed to be multi-rotors, specifically quadrotors, with limited velocity, acceleration, and payload capacity. Each UAV may have a different payload capacity, while the velocity and acceleration constraints remain the same for all UAVs.

\(^1\)https://aerial-core.eu
Ground-based refilling stations are positioned along the power line, providing safe locations where the multi-rotors can load new diverters within a known and bounded time interval to resume their operations. The primary objective is to plan trajectories for the multi-rotors, ensuring that the mission is completed within the specified maximum time. This planning process must account for the vehicle dynamics and capacity constraints, as well as safety requirements such as obstacle avoidance and maintaining a safe distance between UAVs. A map of the environment, which includes a polyhedral representation of obstacles such as power towers and cables, is assumed to be available prior to the installation mission.

III. PRELIMINARIES

Let us consider a discrete-time dynamical model of a multi-rotor system $x_{k+1} = f(x_k, u_k)$, where $x_{k+1}, x_k \in \mathcal{X} \subseteq \mathbb{R}^n$ represent the next and current states of the system, respectively, and $u_k \in \mathcal{U} \subseteq \mathbb{R}^m$ is the control input. Let us also assume that $f : \mathcal{X} \times \mathcal{U} \rightarrow \mathcal{X}$ is differentiable in both arguments and locally Lipschitz. Therefore, given an initial state $x_0 \in \mathcal{X}_0 \subseteq \mathbb{R}^n$ and a time vector $t = (t_0, \ldots, t_N)^T \in \mathbb{R}^{N+1}$, with $N \in \mathbb{N}_0$ being the number of samples that describe the evolution of the system with the sampling period $T_s \in \mathbb{R}_{>0}$, we can define the finite control input sequence $u = (u_0, \ldots, u_{N-1})^T \in \mathbb{R}^N$ as the input to provide the system to obtain the unique sequence of states $x = (x_0, \ldots, x_N)^T \in \mathbb{R}^{N+1}$.

Hence, we can define the state sequence $x$ and the control input sequence $u$ for the $d$-th multi-rotor, with $d \in \mathbb{N}_{>0}$, as $\dot{x} = (d\mathbf{p}(j), d\mathbf{v}(j), d\mathbf{a}(j))^T$ and $d\mathbf{u} = (d\mathbf{a}(1), d\mathbf{a}(2), d\mathbf{a}(3))^T$, where $d\mathbf{a}(j)$, $d\mathbf{v}(j)$, and $d\mathbf{p}(j)$, with $j = \{1, 2, 3\}$, represent the sequences of position, velocity, and acceleration of the vehicle along the $j$-axis of the world frame, respectively. The $k$-th elements of $d\mathbf{p}(j)$, $d\mathbf{v}(j)$, $d\mathbf{a}(j)$, and $t$ are denoted as $d\mathbf{p}_k(j)$, $d\mathbf{v}_k(j)$, $d\mathbf{a}_k(j)$, and $t_k$, respectively, where $k \in \mathbb{N}_{>0}$.

The multi-rotor system model $x_{k+1} = f(x_k, u_k)$ employed in this paper is based on the formulation presented in the authors’ prior work [47]. In this formulation, motion primitives are represented using splines, offering advantages such as generating trajectories with arbitrary maneuver durations and state constraints. The computational burden is manageable, even on standard laptops, handling millions of motion primitives per second. Splines allow independent approximation of the multi-rotor’s translational dynamics along each $j$-axis [50]. Here, we simply refer to them as $\left(d\mathbf{p}_{k+1}(j), d\mathbf{v}_{k+1}(j), d\mathbf{a}_{k+1}(j)\right)^T = d\mathbf{p}(j), d\mathbf{v}(j), d\mathbf{a}(j)$.

Throughout the following discussions, a label will be assigned to indicate the specific drone associated with the dynamic system, while the vector stack containing all drone variables will not include any explicit labels.

A. SIGNAL TEMPORAL LOGIC

Definition 1 (Signal Temporal Logic): STL was introduced for the first time in [17] to monitor the behavior of real-valued signals. STL provides a concise and unambiguous representation of complex system behaviors by encoding all mission specifications within a single formula $\varphi$ [17]. For instance, it can express requirements such as “at least two vehicles need to survey regions A and B, one must visit region C within the time interval $[t_1, t_2]$, while all vehicles must comply with safety requirements”. A comprehensive explanation of the STL syntax and semantics can be found in [15] and [17], but is omitted here for brevity. In short, an STL formula $\varphi$ is defined based on a set of predicates $p_i$, with $i \in \mathbb{N}_0$, which are atomic prepositions representing simple operations such as belonging to a region or comparisons of real values. These predicates can be combined using Boolean operators such as negation ($\neg$), conjunction ($\land$), disjunction ($\lor$), and implication ($\Rightarrow$), as well as temporal operators including eventually ($\Diamond$), always ($\Box$), until ($\mathcal{U}$), and next ($\mathcal{N}$). The resulting STL formula $\varphi$ is considered valid if the expression evaluates to true ($\top$), and invalid ($\bot$) otherwise. For instance, informally, the expression $\varphi_1 \mathcal{U} \varphi_2$ implies that predicate $\varphi_2$ must hold at some point within the time interval $I$, and until then, predicate $\varphi_1$ must remain continuously satisfied.

B. ROBUST SIGNAL TEMPORAL LOGIC

Definition 2 (STL Robustness): The presence of system uncertainties, a dynamic environment, and unforeseen events can affect the satisfaction of an STL formula $\varphi$. To account for a margin of maneuverability in meeting $\varphi$, measuring how well (poorly) a given specification is satisfied, the concept of robust semantics for STL formulae has been proposed [15], [17], [51]. This robustness, denoted as $\rho$, is a quantitative metric that guides the optimization problem derived towards the best feasible solution to achieve mission satisfaction. This value can be formally described by using the following formulae in a recursive manner:

$$\rho_{\varphi}(x, t_k) = \mu(x, t_k),$$
$$\rho_{\neg \varphi}(x, t_k) = -\rho_{\varphi}(x, t_k),$$
$$\rho_{\varphi_1 \vee \varphi_2}(x, t_k) = \min\left(\rho_{\varphi_1}(x, t_k), \rho_{\varphi_2}(x, t_k)\right),$$
$$\rho_{\varphi_1 \wedge \varphi_2}(x, t_k) = \max\left(\rho_{\varphi_1}(x, t_k), \rho_{\varphi_2}(x, t_k)\right).$$
\( \rho_{\Box \varphi}(x, t_k) = \min_{t_i' \in [t_k + I]} \rho_{\varphi}(x, t_i') \),

\( \rho_{igcirc \varphi}(x, t_k) = \max_{t_i' \in [t_k + I]} \rho_{\varphi}(x, t_i') \),

\( \rho_{\Diamond \varphi}(x, t_k) = \rho_{\varphi}(x, t_i') \), with \( t_i' \in [t_k + I] \),

\( \rho_{\varphi}I \rho_{\varphi}2(x, t_k) = \max_{t_i' \in [t_k + I]} \left( \min_{t_i' \in [t_k + I]} \rho_{\varphi}(x, t_i') \right) \),

where \( t_k + I \) represents the Minkowski sum of the scalar \( t_k \) and the time interval \( I \). The formulae described above consist of a set of predicates, denoted as \( p_i \), along with their corresponding real-valued functions \( \mu_i(x, t_k) \). These predicates are considered true if their robustness value is greater than zero, and false otherwise.

The entire formula operates as a logical expression, where it evaluates to false if at least one predicate is false. The evaluation follows the application of logical and temporal operators (such as \( \mu \), \( \Box \), \( \Diamond \), \( \circ \)) from the innermost part to the outermost part of the formula. For instance, an example could involve being inside a target region or outside an obstacle region, with the regions being defined by a specific number of predicates. Further details can be found in [15], [17], and [52]. In this case, we say that \( \rho \) satisfies the STL formula \( \varphi \) at time \( t_k \) if \( \rho_{\varphi}(x, t_k) > 0 \) (in short, denoted as \( x(t_k) \models \varphi \)), and \( \varphi \) violates \( \rho \) if \( \rho_{\varphi}(x, t_k) \leq 0 \). To simplify notation, we will use \( \rho_{\varphi}(x) \) instead of \( \rho_{\varphi}(x, t_k) \) when \( t_k = 0 \).

Thus, we can compute the control inputs \( u \) that maximize the robustness \( \rho_{\varphi}(x) \) over a set of finite state and input sequences \( x \) and \( u \), respectively. An optimal sequence \( u^* \) is considered valid if \( \rho_{\varphi}(x^*) \) is positive, where \( x^* \) and \( u^* \) obey the dynamical system. A higher value of \( \rho_{\varphi}(x^*) \) indicates a more robust behavior of the system against disturbances, enabling the system to tolerate higher levels of disruption without violating the STL specification.

**Definition 3 (Smooth Approximation [42]):** Let us consider \( \lambda \in \mathbb{R}_{>0} \) as a tunable parameter. The smooth approximation of the min and max operators with \( \beta \)-predicate arguments is:

\[
\begin{align*}
\max(\rho_{\varphi_1}, \ldots, \rho_{\varphi_\beta}) & \approx \sum_{i=1}^{\beta} \rho_{\varphi_i} e^{\lambda \rho_{\varphi_i}}, \\
\min(\rho_{\varphi_1}, \ldots, \rho_{\varphi_\beta}) & \approx -\frac{1}{\lambda} \log \left( \sum_{i=1}^{\beta} e^{-\lambda \rho_{\varphi_i}} \right).
\end{align*}
\]

In this paper, we introduce an improved approximation approach compared to our previous work [47]. This new method possesses the properties of being asymptotically complete and smooth everywhere, similar to the well-known Log-Sum-Exponential (LSE) approximation [15]. Furthermore, it guarantees soundness by ensuring that an optimal sequence \( u^* \) with strictly positive smooth robustness (\( \rho_{\varphi^*}(x) > 0 \)) satisfies the specification \( \varphi \), while an optimal sequence \( u^* \) with negative smooth robustness (\( \rho_{\varphi^*}(x) \leq 0 \)) violates it. The property of *asymptotical completeness* indicates that as the parameter \( \lambda \) tends towards infinity, the approximation of the final smooth robustness formula \( \tilde{\rho}_{\varphi^*}(x) \) can be arbitrarily close to the true robustness \( \rho_{\varphi^*}(x) \). By achieving smoothness everywhere, the approximation is infinitely differentiable, enabling the utilization of gradient-based optimization algorithms to find solutions to the problem at hand [42].

**Definition 4 (STL Motion Planner [47]):** Starting from mission specifications encoded as STL formula \( \varphi \), and replacing its robustness \( \rho_{\varphi^*}(x) \) with the smooth approximation \( \tilde{\rho}_{\varphi^*}(x) \) (Def.3), the generation of multi-rotor trajectories can be cast as an optimization problem:

\[
\begin{align*}
\text{maximize} \quad & \tilde{\rho}_{\varphi}(p^{(j)}, v^{(j)}) \\
\text{s.t.} \quad & d_{\varphi}^{(j)} \leq d_{\rho}^{(j)} \leq d_{\varphi}^{(j)} \\
& d_{\varphi}^{(j)} \leq d_{\rho}^{(j)} \leq d_{\varphi}^{(j)} \\
& \tilde{\rho}_{\varphi}(p^{(j)}, v^{(j)}) \geq \varepsilon, \\
& d_{S^{(j)}}^{(j)}, \forall k = \{0, 1, \ldots, N - 1\},
\end{align*}
\]

where \( d_{\varphi}^{(j)} \) and \( d_{\rho}^{(j)} \) represent the desired maximum values for velocity and acceleration, respectively, of drone \( d \) along each \( j \)-axis of the world frame. The set of drones is denoted as \( D \). The lower bound on robustness, \( \tilde{\rho}_{\varphi}(p^{(j)}, v^{(j)}) \geq \varepsilon \), provides a safety margin for satisfying the STL formula \( \varphi \) in the presence of disturbances. As demonstrated in [53], disturbances with magnitudes smaller than \( \varepsilon \) do not lead to a violation of the formula. The value of \( \varepsilon \) can be computed such that \( |\rho_{\varphi}(x) - \tilde{\rho}_{\varphi}(x)| \leq \varepsilon \). Finally, \( d_{S^{(j)}}^{(j)} \) refers to the motion primitives employed to describe the motion of drone \( d \) along each \( j \)-axis, as previously mentioned in the Preliminaries.

**IV. PROBLEM SOLUTION**

In this section, we utilize the STL framework presented in Section III to formulate the optimization problem outlined in Section II. This formulation results in a nonlinear non-convex max-min problem, which we represent as a Nonlinear Programming (NLP) formulation and solve using dynamic programming techniques (Section IV-A). Solving this type of nonlinear problem within a reasonable time frame is challenging due to the propensity of solvers to converge to local optima [48], [49]. To address this issue, we compute an initial guess using a simplified MILP formulation in Section IV-B, which does not consider obstacle avoidance, safety requirements, vehicle dynamics, and time mission specifications. This simplification facilitates the search for a global solution. Additionally, our proposed framework includes a mechanism for mission replanning in the event of UAV failures (Section IV-C). Finally, we enhance the original motion

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2The computation of \( \rho_{\varphi^*}(x) \) involves non-differentiable functions such as \( \min \) and \( \max \). To address the computational complexity associated with these non-differentiable functions, it is beneficial to employ a smooth approximation, denoted as \( \tilde{\rho}_{\varphi^*}(x) \), of the robustness function \( \rho_{\varphi^*}(x) \). This smooth approximation provides a more tractable and computationally efficient alternative.
In this section, we extract the mission specifications for implicit savings in multi-rotor flight time (Section IV-D). We aim to minimize energy consumption, leading to a collaborative team of UAVs that can effectively respond to unforeseen events, cover large-scale scenarios, and minimize operation time. The mission specifications can be categorized into two types: safety requirements and task requirements. The safety requirements ensure the overall safety of the operation throughout the entire mission time $T_N$. These requirements include constraints such as the UAVs staying within the designated workspace $(d\varphi_{ws})$, avoiding collisions with obstacles in the environment $(d\varphi_{obs})$, and maintaining a safe distance from other UAVs $(d\varphi_{dis})$. On the other hand, the task requirements focus on achieving specific tasks at predefined time intervals during the mission. These requirements involve ensuring that all target regions are visited by at least one UAV $(d\varphi_{tr})$ and that each UAV remains in a target region for the designated installation time $T_{ins}$. Considering the limited payload capacity of the UAVs, they are required to visit a refilling station and remain there for a refilling time $T_{rs}$ once they run out of onboard diverters $(d\varphi_{rs})$. Notably, we do not explicitly specify the number or sequence of target regions that a single drone must visit. Rather, we grant the framework the flexibility to determine the most convenient number and sequence of targets for each vehicle, while still allowing for the possibility of visiting all targets if needed. Finally, after completing their installation operations, each UAV should fly to the closest refilling station $(d\varphi_{hm})$. By incorporating all these mission specifications, the STL formula can be formulated as follows:

$$\varphi = \bigwedge_{d \in D} \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{rs}, N_{ins}} (d\varphi_{ws} \land d\varphi_{obs} \land d\varphi_{dis})$$

$$\land \bigvee_{q=1}^{N_{rs}, N_{ins}} \bigwedge_{d \in D} \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{ins}} (d c(t_k) > 0) (d\varphi_{tr,q})$$

$$\land \bigvee_{q=1}^{N_{rs}, N_{ins}} \bigwedge_{d \in D} \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{rs}} \bigwedge_{d \in D} \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{ins}} \bigwedge_{d \in D} \left( (d c(t_k) = 0 \Rightarrow (d\varphi_{ins}) \land (d\varphi_{ws})) \land (d\varphi_{obs}) \land (d\varphi_{dis}) \land (d\varphi_{hm}) \right).$$

(2)

The STL formula $\varphi$ consists of six specifications $(d\varphi_{ws}, d\varphi_{obs}, d\varphi_{dis}, d\varphi_{tr}, d\varphi_{rs}, d\varphi_{hm})$ and three time intervals $(T_N, T_{ins}, T_{rs})$. The following equations describe each of these specifications:

$$d\varphi_{ws} = \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{ins}} p^{(j)}(t_k) \in (p^{(j)}_{ws}, \tilde{p}^{(j)}_{ws}),$$

(3a)

$$d\varphi_{obs} = \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{ins}} p^{(j)} \notin (p^{(j)}_{obs}, \tilde{p}^{(j)}_{obs}),$$

(3b)

$$d\varphi_{dis} = \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{ins}} \bigwedge_{d \in D} \bigwedge_{m \neq d} \| d p - m p \| \geq \Gamma_{dis},$$

(3c)

$$d\varphi_{hm} = \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{ins}} (p^{(j)}_{hm} \in (p^{(j)}_{hm}, \tilde{p}^{(j)}_{hm})),$$

(3d)

$$d\varphi_{tr,q} = \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{ins}} (p^{(j)} \in (p^{(j)}_{tr,q} \land \tilde{p}^{(j)}_{tr,q})) \land$$

$$\land \bigwedge_{q=1}^{N_{ins}} (d c(t_k) = 0) (d c(t_k) - 1) - 1),$$

(3e)

$$d\varphi_{rs,q} = \bigwedge_{j=1}^{N_{tr}} \bigwedge_{q=1}^{N_{ins}} (p^{(j)} \notin (p^{(j)}_{rs,q} \land \tilde{p}^{(j)}_{rs,q}) \land$$

$$\land \bigwedge_{q=1}^{N_{ins}} (d c(t_k) = 0) (d c(t_k) + 1),$$

(3f)

where (3a) constrains the position of each UAV to remain within the designated workspace, with $p^{(j)}_{ws}$ and $\tilde{p}^{(j)}_{ws}$ representing the working space limits. Obstacle avoidance and mission completion operation are defined in (3b) and (3d), respectively. Rectangular regions with vertices denoted by $p^{(j)}_{obs,q}$, $p^{(j)}_{hm}$, $p^{(j)}_{obs}$, and $\tilde{p}^{(j)}_{hm}$ define obstacle and drone’s final position (closest refilling station) areas. The safety distance requirement is encoded in (3c), where $\Gamma_{dis} \in \mathbb{R}_{>0}$ represents the threshold value for the mutual distance between UAVs. Finally, visiting target areas and refilling stations are described by (3e) and (3f), respectively. The vertices $p^{(j)}_{tr,q}$, $p^{(j)}_{rs,q}$, $\tilde{p}^{(j)}_{tr,q}$, and $\tilde{p}^{(j)}_{rs,q}$ identify the target and refilling areas. In (3e) and (3f), the payload capacity $(d c(t_k)) \in \mathbb{N}_0$ is used to represent the number of remaining and loaded diverters for each UAV, respectively. The payload capacity is constrained to the interval $0 \leq (d c(t_k)) \leq (d c_{max})$, where $(d c_{max})$ corresponds to the nominal payload capacity. The always operators $(\square)$ guarantee the fulfillment of the time requirements $T_{ins}$ and $T_{rs}$, which correspond to the durations of the installation and refilling operations, respectively. On the other hand, the next operators $(\bigcirc)$ ensure that changes in the payload capacity will occur only following those time intervals. Additionally, the eventually operators $(\bigtriangledown)$ ensure that the mission requirements will ultimately be fulfilled within the designated mission time $T_N$.

Using the specifications described in (2), the optimization problem defined in Def.4 is formulated to compute feasible trajectories that maximize the smooth robustness $\tilde{p}_q(x)$ with respect to (w.r.t) the given mission specifications $\varphi$. To achieve this, it is necessary to compute the robustness score for each predicate. The STL formula (2) comprises two types of predicates. The first type evaluates whether the UAV position belongs or does not belong to a specific region, as depicted in (3a), (3b), (3d), (3e), and (3f). The second type evaluates the distance between UAVs, as represented by the safety requirement in (3c). The robustness values are quantified based on the Euclidean distance. For predicates belonging to the first group, a positive robustness indicates that the UAV lies within the designated region. The robustness increases as the minimum Euclidean distance to the boundaries of the region along the trajectory becomes larger. However, in the case of (3b), the inverse applies, where being within the obstacle region corresponds to a negative robustness. In the safety distance predicate (3c), the robustness is positive when the distance between UAVs exceeds the threshold $\Gamma_{dis}$. The robustness value increases as the minimum Euclidean distance between UAVs along the
In other words, we can express this as

\[ \bar{d} \rho = \min_{p} \left( \min(d\rho_{g_{ij}}^{(1)}, d\rho_{g_{ij}}^{(3)}, d\rho_{g_{ij}}^{(2)}, d\rho_{g_{ij}}^{(4)}) \right), \]

where

\[ d\rho_{g_{ij}}^{(k)} = \bar{p}_{ws}^{(k)} - d_{k}^{(k)} = d_{k}^{(k)} - p_{ws}^{(k)}. \]

Similarly, the robustness score of the non-belonging predicate \( d\rho_{obs} \) (3b) can be computed by taking the inverse of each minimum distance for each sample along the trajectory. In other words, we can express this as:

\[ d\rho_{obs} = \min_{p} \left( -\min(d\rho_{g_{ij}}^{(1)}, d\rho_{g_{ij}}^{(3)}, d\rho_{g_{ij}}^{(2)}, d\rho_{g_{ij}}^{(4)}) \right), \]

where \( d\rho_{g_{ij}}^{(k)} \) and \( d\rho_{g_{ij}}^{(k)} \), with \( j = \{1, 2, 3\} \), can be computed by replacing \( p_{ws}^{(k)} \) and \( p_{obs}^{(k)} \) in (4) with \( p_{obs}^{(k)} \) and \( p_{obs}^{(k)} \), respectively. On the contrary, the robustness score of the safety requirement predicate \( d\rho_{dis} \) (3c) can be expressed as:

\[ d\rho_{dis} = \min_{p} \left( \sum_{d, d \neq m} \left( \| d - m \| - \Gamma_{dis} \right) \right), \]

where the minimum distance between UAVs is computed for each time step \( t_k \) in the trajectory, and then the min value of that vector of minimum distances.

It is important to emphasize that the interdependence of mission requirements necessitates the joint planning of all UAV trajectories through a unified optimization problem.

**B. INITIAL GUESS**

To ensure optimal solutions within a reasonable time and prevent the solver from becoming trapped, it is essential to have an appropriate initial guess for the STL motion planner. The key concept behind obtaining this initial guess is to simplify the original diverter installation problem by abstracting it into an optimization with fewer constraints. Specifically, the initial guess focuses on fulfilling mission requirements related to visiting target regions, considering refilling and mission completion operations \( d\rho_{gs}, d\rho_{ev}, \) and \( d\rho_{hm} \). It omits obstacle avoidance, workspace, and safety distance requirements \( d\rho_{obs}, d\rho_{ws}, \) and \( d\rho_{dis} \), and the mission time intervals \( (T_k, T_{max}, \) and \( T_{dis} \)) since these constraints introduce complex nonlinearities and motion discontinuities in the problem.

To generate the initial guess, a graph-based representation is employed, establishing connections between target regions and refilling stations. This formulation is modeled as a variant of the Vehicle Routing Problem (VRP) using Mixed-Integer Linear Programming. Assigning target regions to vehicles and providing navigation sequences for each UAV are the objectives of the MILP solution.

Figure 3 illustrates the graph used in our approach, which is an undirected weighted multigraph denoted by the tuple \( G = (V, E, W, D, \mathcal{C}) \). The set of vertices, \( V \), consists of locations that encompass target regions, refilling stations, and depots where each UAV is initially positioned at time \( t_0 \). Specifically, \( V = T \cup R \cup O \), with \( |T| = \tau \) representing the target regions, \( |R| = r \) representing the refilling stations, and \( |O| = \delta \) including the depots. The set of edges, \( E \), and their corresponding weights, \( W \), define the connectivity and distances between the vertices, respectively. Additionally, as before, \( D \) represents the set of UAVs, with \( |D| = \delta \), and \( \mathcal{C} = \{1^\circ, \ldots, \delta^\circ\} \) denotes their respective maximum payload capacities.

Regarding the graph connectivity, all vertices in \( T \) are fully connected to each other and connected to every vertex in the set \( R \cup O \). Each connection involves one edge per UAV, totaling \( \delta \) edges. More formally, let \( e_{ijd} \in E \) be the edge that connects vertices \( i \) and \( j \), with \( i, j \in V \) and \( i \neq j \), using UAV \( d \in D \), where \( d = \{1, \ldots, \delta\} \). The weight associated with \( e_{ijd} \) is denoted as \( w_{ijd} \in W \). Considering the homogeneous dynamic constraints for UAVs, such as maximum velocities \( \bar{\rho} \) and maximum accelerations \( \bar{\rho} \), and assuming similar flight durations for each UAV, the mission operation constraint can be expressed as a maximum distance traveled by each individual UAV. As a result, we model the edge weights using Euclidean distances, resulting in \( w_{ijd} = d_{ijd} \).

For each edge \( e_{ijd} \in E \), we introduce an integer variable \( z_{ijd} \in \mathbb{Z}_0 \). This variable represents the number of times the corresponding edge is selected in the solution of the MILP. Thus, for a given UAV \( d \in D \), we have \( z_{ijd} \in \{0, 1\} \) for all \( \{i, j\} \in \{T, O\} \), and \( z_{ijd} \in \{0, 1, 2\} \) if \( i \in R \) and \( j \in T \). The former ensures that the edge between two target regions is never covered twice, indicating that the UAV starts from a depot and never returns to it. The latter allows for round trips between refilling stations and target regions, in case there are no other diverters to be installed. Specifically, \( z_{ijd} = 2 \) denotes a return trip between a refilling station \( i \) and a target region \( j \) of drone \( d \). When the script of the variable \( z_{ijd} = 0 \), it indicates that the corresponding vertex \( i \) or \( j \) is a depot, depending on the order. Utilizing these variables, we can formulate the MILP problem as follows:

\[
\text{minimize} \quad \sum_{\{i, j\} \in V, i \neq j, d \in D} w_{ijd} z_{ijd} \tag{7a}
\]
\[
\text{s.t.} \quad \sum_{i \in V, \not i \neq j, d \in D} z_{ijd} = 2, \forall j \in T. \tag{7b}
\]
In the above formulae, the objective function (7a) quantifies the total distance covered by the multi-UV team. Constraints (7b) and (7c) ensure that each target region is visited exactly once. To achieve this, auxiliary integer variables $y_{jd}$ $\in \{0, 1\}$ are introduced, which ensure that if a UAV $d \in D$ reaches target $j \in T$, the same UAV must also leave the target. Constraint (7d) guarantees that each UAV starts the mission from its depot and does not return to it. Constraints (7e) serve two purposes: preventing tours that exceed the payload capacity of the UAVs and ensuring that all tours connect to a refilling station, which is commonly known as the sub-tour elimination constraint [54], [55]. The lower bound $h(T)$ represents the minimum number of UAVs required to visit all target regions $T$. In practice, $h(T)$ is often defined as $\left\lceil \sum_{i \in T} c_i/C \right\rceil$, where $c_i$ is the number of diverters at target $i$ and $C$ is the payload capacity of the UAVs. To address the computational challenges posed by the exponential number of constraints associated with $h(T)$ [56], [57], [58], an adaptive constraint incorporation strategy is employed. This approach dynamically adds constraints as violations occur during the optimization process [55]. By incorporating constraints only when necessary, based on violations identified during the solution process, the computational burden is alleviated. This adaptive constraint management strategy improves the computational efficiency of the solution by avoiding the exhaustive evaluation of all possible combinations and focusing on constraints that are most relevant or impactful. After solving the MILP problem, the optimal assignment of target regions and refilling stations for each UAV is obtained. To generate dynamically feasible trajectories, we employ multi-rotor motion primitives [47], as said in the Preliminaries (Section III). The detailed computation of these trajectories is provided in [50] and is not presented here for brevity. In summary, the motion primitives are computed by assuming rest-to-rest motion between regions, where the UAV has zero velocity and acceleration. The minimum feasible time required to satisfy the desired maximum values of velocity $\dot{v}^{(j)}$ and acceleration $\ddot{a}^{(j)}$ along the UAV’s trajectory is imposed. The trajectories also account for the installation time intervals $T_{ins}$ and refilling time intervals $T_{rs}$ during which the UAVs remain stationary at the corresponding regions.

**C. EVENT-TRIGGERED REPLANNER**

The presented motion planner enables the generation of feasible trajectories for a multi-UAV team carrying out a bird diverter installation mission. However, in real-world scenarios, UAV failures can occur (e.g., due to a battery or technical fault), resulting in the need for an alternative UAV to take over the pending tasks to ensure the successful completion of the mission. To address this, an online event-based replanning procedure is employed.

Whenever a UAV failure event is detected, a new plan is computed online using the motion planner described in Sections IV-A and IV-B. The trajectories of the UAVs that are still operational are not recomputed since they can continue with their original plans. The replanning process focuses on generating a new trajectory for the backup UAV that replaces the faulty one. Assuming the failure event occurs at time $t_{fail} \in [t_0, T_N]$, the new trajectory for the backup UAV is generated within the time interval $[t_{fail} + T_{rep}, T_N]$. Here, $T_{rep}$ represents the maximum expected computation time for replanning, which was estimated by running multiple instances of the STL optimization problem and varying factors such as the number and positions of target regions, as well as the time instance $t_{fail}$ when the failure occurs.

During the replanning process, the motion planner utilizes the original trajectories of the non-faulty UAVs within the time interval $[t_{fail} + T_{rep}, T_N]$ as input constraints for the safety distance requirement ($d_{\text{dis}}$). As said before, these trajectories do not need to be recomputed. Additionally, only the pending target regions that were previously assigned to the faulty UAV are considered in the replanning procedure. However, the order in which these target regions are visited may change, as the backup UAV may start from a different position or with a different payload, potentially resulting in a suboptimal visitation sequence.

It is important to note that the computational effort required for replanning is significantly lower than that for initial planning, as it involves only a single UAV and a reduced set of target regions (those that have not yet been visited). While exploring alternative re-planning strategies for the routes of non-faulty drones may potentially yield more optimal solutions in terms of robustness scores, it is essential to emphasize that our paper focuses on ensuring operational continuity, minimizing disruptions, and maintaining safety in hazardous scenarios. Prioritizing the better trajectories solution in these cases may lead to the drones halting their mission, resulting in delays and increased costs for the entire operation. Therefore, our approach offers a rapid and effective solution by replacing a faulty multi-rotor with a backup vehicle, ensuring seamless continuation of the mission.

**D. ENERGY-AWARE PLANNER**

This section introduces an energy-aware extension of the motion planner presented in Section IV-A, focusing on minimizing the energy consumption of the UAVs during the mission. This results in implicit savings in multi-rotor flight time. The key concept is to determine the optimal cruising speed, considering the aerodynamics of the multi-rotors, to achieve minimal energy usage. The goal is to generate trajectories for the UAVs that closely align with these optimal speeds.
The power required by a multi-rotor decreases as the low forward speed increases, while it increases significantly at high speeds due to parasitic losses. This combination of effects gives rise to an optimal forward speed that minimizes the energy consumption of the UAV and maximizes its flight time. This energetically optimal forward speed, denoted as \( v^* \), is described in [59, Ch. 5], and defined as:

\[
v^* = \sqrt{\frac{m}{2ρ}} \left( \frac{4kA}{n_f} \right)^{1/4},
\]

where \( m \in \mathbb{R} \) represents the mass of the vehicle, \( n_r \in \mathbb{N} \) denotes the number of rotors, \( A \in \mathbb{R} \) corresponds to the rotor disk area, \( f \in \mathbb{R} \) represents the equivalent flat plate area of the fuselage, \( κ \in \mathbb{R} \) is the induced power correction factor, and \( ρ \in \mathbb{R} \) denotes the air density.

Hence, the optimization problem (1) can be reformulated by including an energy term to prioritize solutions where the forward speed along the trajectory is as close as possible to the optimal value \( v^* \) in terms of energy consumption. Namely, we write:

\[
\begin{aligned}
& \text{maximize} & & \tilde{ρ}_v(p^{(d)}, v^{(d)}) - η \sum_k \left( 1 - \frac{d_{v\text{for}}(d, v_k)}{v^*} \right)^2 \\
& \text{s.t.} & & \frac{d_k}{v^*} \leq \frac{d_{v_k}}{v^*} \leq \frac{d_k}{v^*}, \\
& & & \frac{d_{\text{for}}(d, v_k)}{v^*} \geq \epsilon, \\
& & & \forall k \in \{0, 1, \ldots, N - 1\},
\end{aligned}
\]

where \( η \in [0, 1] \in \mathbb{R} \) serves as a tunable weight for the energy term in the objective function, which quantifies the deviation of the forward speed \( d_{v\text{for}}(d, v_k) \) from the optimal value \( v^* \). The forward speed \( d_{v\text{for}}(d, v_k) \) is computed as the Euclidean norm of the velocity components of drone \( d \) along the \( x \) and \( y \) axes:

\[
d_{v\text{for}}(d, v_k) = \sqrt{\left( \frac{d_{v_k}(1)}{v^*} \right)^2 + \left( \frac{d_{v_k}(2)}{v^*} \right)^2}.
\]

Note that, as demonstrated in [60] and [61], the energy expended by multi-rotors during ascent, descent, and hovering maneuvers is constant since these operations are performed at a constant speed. Therefore, in the context of the optimization problem (9), we can simplify the computation by excluding the \( d_{v(3)} \) component from the forward speed calculation in (10), reducing the computational burden.

It is important to consider the scenario where \( v^* \) exceeds the maximum feasible vehicle speed \( d_{v(0)} \). In such cases, we can reformulate the optimization problem (9) by replacing \( v^* \) with \( \|d_{v(1)} + d_{v(2)}\| \), which represents the feasible forward speed closest to \( v^* \) as defined in (8).

The introduction of the new energy term in the objective function represents a quadratic error that measures the deviation of the vehicle speed along the trajectory from the speed that minimizes energy consumption. It is worth noting that the solutions generated by the energy-aware planner still yield feasible trajectories that satisfy all mission requirements. However, these trajectories strike a balance between energy consumption and potential trade-offs in smooth robustness score \( \tilde{ρ}_v(x) \).

V. EXPERIMENTAL RESULTS

To validate and assess the effectiveness of our proposed planning approach, we conducted a series of simulations and experiments. Initially, numerical simulations were carried out using MATLAB, where the actual vehicle dynamics and trajectory tracking controller were not explicitly modeled. This allowed us to evaluate the planning algorithm’s performance and gain valuable insights into its behavior. Subsequently, to further verify the feasibility of the generated trajectories and leverage the benefits of software-in-the-loop simulations [62], [63], we conducted additional simulations using the Gazebo robotics simulator. Finally, field experiments performed in a mock-up scenario that closely resembles real-world conditions demonstrated the practical applicability of the proposed method.

These experiments aimed to showcase several aspects: (i) the adherence of the planned trajectories to the mission requirements, (ii) the necessity of the STL motion planner in meeting the mission specifications, as the MILP formulation alone is insufficient, (iii) the capability of our approach to handle UAV failures and perform mission replanning, (iv) a comparison of the solutions obtained with and without the energy-aware feature, and (v) the feasibility of our method in real-world scenarios.

The optimization algorithm was implemented using MATLAB R2019b, with the MILP formulated using the CVX framework,\(^3\) and the STL motion planner implemented using the CasADi library\(^4\) with IPOPT\(^5\) as the solver. Within CVX, the choice of heuristic for solving the MILP problem was left to the framework itself. All numerical simulations and experiments were conducted on a computer equipped with an 17-8565U processor (1.80 GHz) and 32GB of RAM, running on the Ubuntu 20.04 operating system. For more detailed information and visual demonstrations of the experimental results, we provide videos that can be accessed at http://mrs.felk.cvut.cz/bird-diverter-ar.

A. BIRD DIVERTERS INSTALLATION

The proposed planning strategy was tested in a bird diverter installation scenario with two UAVs operating in a mock-up scenario (50 m × 20 m × 15 m). The scenario had seven target regions and four refilling stations, as shown in Figure 2. The parameters used for the optimization problem are listed in Table 1. To avoid long waiting periods, symbolic values were used for the installation time \( T_{\text{ins}} \) and refilling time \( T_{\text{ref}} \). Heading angles were adjusted to align the UAVs with the displacement direction or the cables during installation. It is assumed that an onboard low-level controller, e.g. [8], [10],

\(^3\)http://cvxr.com/cvx/
\(^4\)https://web.casadi.org/
\(^5\)https://coin-or.github.io/Ipopt/
TABLE 1. Parameter values for the optimization problem.

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Symbol</th>
<th>Value</th>
<th>Parameter</th>
<th>Symbol</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Payload capacity (UAV1)</td>
<td>$P_1$</td>
<td>2</td>
<td>Payload capacity (UAV2)</td>
<td>$P_2$</td>
<td>4</td>
</tr>
<tr>
<td>STL safety margin</td>
<td>$\gamma$</td>
<td>0.2</td>
<td>Safety initial distance</td>
<td>$\Gamma_{\text{in}}$</td>
<td>3 m</td>
</tr>
<tr>
<td>Max. velocity</td>
<td>$v_{\text{max}}$</td>
<td>3.1 [m/s]</td>
<td>Max. acceleration</td>
<td>$a_{\text{max}}$</td>
<td>3.1 [m/s²]</td>
</tr>
<tr>
<td>Optimal forward speed</td>
<td>$v_{\text{opt}}$</td>
<td>2.5 [m/s]</td>
<td>Time replanning</td>
<td>$T_{\text{r}}$</td>
<td>10 s</td>
</tr>
<tr>
<td>Mission time</td>
<td>$T_m$</td>
<td>15 s</td>
<td>Installation time</td>
<td>$T_{\text{i}}$</td>
<td>5 s</td>
</tr>
<tr>
<td>Refilling time</td>
<td>$T_{\text{r}}$</td>
<td>12 s</td>
<td>Sampling period</td>
<td>$T_{\text{s}}$</td>
<td>0.05 s</td>
</tr>
<tr>
<td>Number of samples</td>
<td>$N$</td>
<td>1360</td>
<td>Scaling factor</td>
<td>$\lambda$</td>
<td>10</td>
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<tr>
<td>Weight step size</td>
<td>$w$</td>
<td>5.1</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Figure 6. Position, velocity, acceleration, safety mutual distance, and payload capacity for “UAV1” and “UAV2” from real-world experiments. The installation and refilling time windows are highlighted in blue and green, respectively.

B. COMPARISON WITH THE INITIAL GUESS SOLUTION

As outlined in Section IV, the MILP formulation serves as an initial seed for the STL optimization problem, enabling the generation of feasible trajectories for the bird diverter installation mission. While the MILP solution is computationally efficient, it is insufficient on its own due to the omission of critical nonlinear aspects such as obstacle avoidance, safety distance requirements, and time constraints. It is important to note that the STL optimization problem, when not seeded with the MILP initial guess, does not converge to a feasible solution. Therefore, a direct comparison between the STL (1) and the MILP (7) solutions cannot be made. Nonetheless, we evaluate the compliance of both solutions with the mission requirements and assess how the proposed hierarchical planner addresses nonlinear complexities such as obstacle avoidance, safety distance, and time requirements, which the MILP alone cannot handle. Furthermore, we demonstrate how the subsequent STL optimization refines the initial solution obtained from the MILP formulation, striving for maximum robustness and satisfying all mission specifications. Figure 7 illustrates the dynamically feasible trajectories obtained using multi-rotor motion primitives [47], along with the sequences of waypoints assigned to each UAV by the MILP solution.
Upon initial examination, it is evident that the MILP formulation does not explicitly consider time, resulting in trajectories that may exceed the specified time limit $T_N$ and become impractical for real-world applications. This limitation also impacts the trajectory quality, as the MILP approach does not control the vehicles’ motion through accelerations $a_j$, but rather sets waypoint sequences for the UAVs to follow, leading to sharp corners and spikes that are challenging to execute in practice. Such abrupt changes in direction can strain the actuators and compromise trajectory tracking, potentially causing deviations from mission specifications and safety requirements, including the risk of collisions with power towers or cables. Additionally, this approach can result in high energy consumption, adversely affecting the mission objectives. Another crucial distinction lies in the treatment of the mutual safety distance constraint ($d_{\phi_{\text{dis}}}$). While the STL optimization problem incorporates this constraint by regulating velocities and accelerations, and thus the vehicles’ positions, without compromising the optimal region visitation sequence, the MILP formulation does not account for vehicle dynamics to mitigate computational complexity. Consequently, adjusting the value of $\Gamma_{\text{dis}}$ (3c) may require entirely different optimal sequences (outputs of the MILP solver) to obtain a feasible solution for the problem.

Figure 8 provides a visual comparison between the solutions obtained by solely solving the MILP formulation and the complete STL optimization problem. The figure clearly illustrates how the STL formulation refines the initial guess solution from the MILP, resulting in trajectories that satisfy obstacle avoidance, safety distance requirements, and other mission specifications. Notably, the trajectories are significantly rearranged to achieve higher robustness values $\tilde{\rho}_\phi(x)$.

Moreover, it is important to highlight that the differences between the STL and MILP solutions extend beyond the trajectory shape and also encompass the sequence of target regions to visit. Figure 9 demonstrates this distinction in a simple scenario involving two UAVs (located at “IP1” and “IP2”) with identical dynamic constraints. The objective is for the UAVs to visit a set of target regions (“TR1”, “TR2”, “TR3”, and “TR4”) within the time interval [0, $T_N$], while adhering to workspace constraints ($d_{\phi_{ws}}$), avoiding obstacles ($d_{\phi_{\text{obs}}}$), and maintaining a safe distance from each other ($d_{\phi_{\text{dis}}}$). In Figure 9(b), the initial guess solution from the MILP formulation assigns target regions based on workload balancing and minimizing travel distance, without considering obstacle avoidance. In contrast, the complete STL planner (Figure 9(a)) reassigns the targets to satisfy all mission specifications. This example highlights the independence of the final STL optimization from the MILP initial guess solution, providing added flexibility to the hierarchical planner. However, it is essential to note that as the problem complexity increases, a significant deviation between the initial guess and the global solution may hinder the STL optimizer, resulting in convergence to local optima.

C. ENERGY-AWARE AND REPLANNING STRATEGY

This section demonstrates the effectiveness of the energy-aware planner and event-triggered replanner, discussed in Sections IV-D and IV-C, respectively, in reducing energy consumption by the UAVs and ensuring mission continuity in the event of UAV failures. The results of numerical simulations conducted in MATLAB are presented in Figures 10 and 12.

Regarding the energy-aware planner, the trajectories incorporating the forward speed term $v^*$ (Figure 10) exhibit longer paths and abrupt changes in direction compared to the trajectories without energy requirements (Figures 4 and 5). This is due to the UAVs needing to approach the energetically optimal forward speed $v^*$ while satisfying the mission...
specifications $\phi$ (eq. (2)) during the installation of diverters. As a result, the optimization problem (9) restricts the range of allowable velocities more than the standard planner (1), where velocity values can be adjusted over a wider range to maximize robustness. Figure 11 provides a comparison of the forward speed $d_{\text{for}}(t, y_k)$ and normalized smooth robustness values $\tilde{\rho}(x)$ for the standard and energy-aware planners. The plot illustrates that higher velocities are associated with a slight decrease in smooth robustness values.

To demonstrate the event-triggered replanner, we conducted a scenario with a pair of UAVs involved in the installation of bird diverters. Initially, both UAVs followed the trajectories shown in Figure 4. Subsequently, we simulated a failure event when “UAV2” (Figure 12) visited a refilling station to replenish diverters. This triggered the execution of the optimization problem (1) and the corresponding MILP (7) for the backup UAV, “UAV3”, stationed at the refilling station “RSA”. The optimization problem considered the remaining unvisited regions from UAV’s mission, “UAV2”, and the trajectory of the non-faulty UAV, “UAV1”. The resulting trajectories are depicted in Figure 12. In this particular scenario, the order of visiting the target regions remained the same, with slight variations observed in the trajectory connecting “TR5” and “RSC” (Figure 4). These trajectory changes were driven by the optimization problem to ensure compliance with safety and obstacle avoidance requirements. The optimization process took approximately 2 s to solve the entire problem, while the MILP solver took only a few
FIGURE 13. System Architecture. The STL Motion Planner on the ground station generates trajectories ($\dot{x}^t$, $\dot{u}^t$, $\dot{\psi}^t$, $\dot{\theta}^t$, $\dot{\phi}^t$) and heading angles ($\psi^t$, $\theta^t$, $\phi^t$) for the multi-rotors. These trajectories and heading angles are used as inputs to the Tracking Controller, which computes thrust ($\tau_T^t$, $\tau_d^t$) and angular velocities ($\omega_d^t$, $\omega_d^t$) for the UAV Plant [47].

tenths of a second. The failure event was detected at $T_{fail} = 9$ s, considering a maximum expected computation time of $T_{rep} = 10$ s.

D. FIELD EXPERIMENTS
Experiments were conducted to evaluate the practical application of the proposed motion planning approach using DJI F450 quadrotors [64], [65] in a mock-up scenario. The scenario, shown in Figure 4, involved two UAVs operating in a workspace that included power towers, cables, and refilling stations. To facilitate experimentation and consider the safety-critical nature of the power line environment, we simulated these objects in MATLAB and Gazebo and validated the trajectories in Gazebo. Videos showcasing simulated and experimental results can be accessed at http://mrs.felk.cvut.cz/bird-diverter-ar.

The system architecture, depicted in Figure 13, incorporates the STL motion planner, which solves the optimization problem (1) to generate trajectories ($\dot{x}^t$, $\dot{u}^t$, $\dot{\psi}^t$, $\dot{\theta}^t$, $\dot{\phi}^t$) for the multi-rotors. The trajectory generation process is performed as a one-shot computation at time $t_0$, and the resulting trajectories serve as references for the UAV trajectory tracking controller [62]. The UAVs were equipped with an Intel NUC computer (i7-8559U processor with 16GB of RAM) and the Pixhawk flight controller. The software stack utilized the Noetic Ninjemys release of ROS running on Ubuntu 20.04. Further details can be found in [64] and [65].

During the real flight experiments, we were able to verify the successful completion of the installation mission defined by the STL formula (2). The flights also demonstrated compliance with physical constraints and safety requirements, including maximum velocity ($\dot{v}^t$), maximum acceleration ($\ddot{a}^t$), and the minimum safety distance ($\Gamma_{dis}^t$), as well as the payload capacity of the vehicles, $\bar{c}$ (Table 1). Figure 14 shows snapshots of the experiments, where the proximity of the UAVs to the mechanical infrastructure (cables and towers) and refilling locations can be observed. To provide this visual representation, we used a ROS package to project 3D mesh files, which were originally used in MATLAB and Gazebo to represent the scenario, onto the camera frames of the UAVs.

VI. DISCUSSION
In this discussion section, we aim to provide a comprehensive analysis and insights into our proposed planning approach in the context of installing bird diverters. The work presented here forms a part of the AERIAL-CORE European project. While we acknowledge the existence of established approaches in the literature, such as the kinodynamic RRT* [20] and the Capacitated Vehicle Routing Problem (CVRP) formulation [66], which were not originally designed for multi-robot systems and do not explicitly consider time requirements like STL, it is important to clarify that our objective is not to challenge or claim superiority over these approaches. Rather, our focus is on leveraging TL techniques and applying them to a real-world use-case scenario.

One of the distinct advantages of using TLs is its ability to capture complex temporal constraints and requirements in a formal and expressive manner. By combining natural language commands, temporal and Boolean operators, and task and motion planning, we can generate trajectories that fulfill mission objectives while considering the dynamics and constraints of the multi-rotor system. While the robustness values can indeed be computed using functions of the variables, the use of TLs provides a more concise and intuitive framework for specifying and addressing these mission specifications.

In our work, we make a deliberate choice to relax certain constraints in the MILP formulation, including obstacle avoidance, workspace limits, safety distance requirements, and mission time intervals (Section IV-B). This decision was made based on several considerations. Firstly, including these constraints would introduce complex nonlinearities and motion discontinuities, making the MILP formulation more challenging and computationally expensive and leading to a non-convex optimization problem. Secondly, considering time constraints would require incorporating motion primitives, further complicating the problem formulation. Adding excessive complexity in this aspect would not be beneficial, but rather counterproductive, as our goal is to obtain solutions as quickly as possible. The MILP formulation serves as a foundation for the subsequent STL optimization problem, acting as a stepping stone for the STL motion planner. Its primary purpose is to enable the avoidance of local optima and facilitate the refinement of the trajectory generation process.

To highlight the significance of our approach, we compare the solutions obtained from the MILP formulation and the complete STL optimization problem. The results, as showcased in Figure 8, demonstrate the refinement achieved by the STL formulation, particularly in addressing mission requirements such as obstacle avoidance and safety distance. This comparison serves to illustrate the added value and effectiveness of our proposed solution in the context of bird diverter installation. Furthermore, we highlight in Section V-B and visually depict in Figure 9 a simple scenario example where we explicitly demonstrate that the differences...
between the STL and MILP solutions extend beyond the trajectory shape, encompassing the sequence of target regions to visit.

Regarding the replanning of other vehicles’ routes (Section IV-C), our focus is on providing a practical solution that ensures continuity of operations in hazardous scenarios where swift action is crucial. We describe an online event-based replanning procedure where a new plan is computed for the backup multi-rotor in the event of a failure, while the trajectories of operational multi-rotors remain unchanged to ensure continuity. The replanning process aims to generate a new trajectory for the backup multi-rotor within a specified time interval, considering factors such as expected computation time. While the evaluation of alternative re-planning strategies for other vehicles’ routes could be an interesting avenue for future research, it is important to note that our paper specifically addresses the installation of bird diverters and prioritizes the practicality and continuity of operations in this specific context. Although it may not provide the optimal solution for all vehicles’ routes, it offers a fast and effective means to replace a faulty multi-rotor with a backup vehicle.

It is important to note that our work is specifically tailored to the practical context of installing bird diverters, which introduces unique challenges and constraints. These characteristics make our work novel and unique, as we have incorporated TL techniques into a problem that has not been extensively explored from this perspective. While existing procedures may address similar problems, our focus on applying TL techniques to this specific real-world scenario adds a valuable contribution to the field.

### A. CHALLENGES OF MULTI-ROBOT BIRD DIVERSION

In this paper, we have explored the advantages and potential of using a multi-rotors team for bird diversion. However, it is important to address the drawbacks and considerations associated with this approach.

Coordination among multiple multi-rotors can be a complex task, requiring sophisticated control algorithms and communication protocols to ensure proper synchronization and cooperation. Effective communication is crucial for coordination and information sharing, but it can be susceptible to interference, signal loss, or limited bandwidth, which may disrupt the exchange of critical data. Furthermore, the presence of multi-rotors in the vicinity of birds may introduce disturbances that can affect their behavior. The noise and visual presence of the vehicles can cause stress or alter the natural patterns of bird movement. It is crucial to assess and mitigate these potential disturbances to minimize adverse effects on the target birds and their habitats.

Cost is another important consideration. The scalability and economic feasibility of employing a multi-rotors team for bird diversion, especially considering the large number of power towers, need to be carefully evaluated. Further research is required to optimize resource allocation and operational strategies, aiming to minimize costs while maintaining effectiveness. Moreover, the impact of wind and extreme weather on drone trajectories should be taken into account. Strong winds can affect the stability and maneuverability of multi-rotor systems, potentially hindering their effectiveness in bird diversion tasks. Integrating weather forecasting and real-time
adaptation mechanisms can enhance safety and performance under varying environmental conditions.

VII. CONCLUSION
This paper has presented a motion planning framework for encoding bird diverter installation missions for a team of multi-rotors with payload capacity limitations and dynamic constraints. The proposed method utilizes STL specifications to generate dynamically feasible trajectories that satisfy mission requirements, including safety and mission time constraints. The work builds upon a previous motion planner by adopting an MILP approach to handle the nonlinear non-convex optimization problem. The MILP solver provides an initial guess solution for the STL framework, facilitating algorithm convergence. Additionally, event-triggered replanning and energy-aware planning techniques are introduced to address UAV failures and optimize energy consumption, respectively. MATLAB and Gazebo simulations, along with field experiments, validate the effectiveness of the proposed approach in a real-world scenario.

Our analysis reveals that the simplified MILP solver alone is insufficient for the application, but it serves as a valuable initial guess for the complete STL planner. The hierarchical approach allows us to address a broader range of mission specifications and requirements compared to existing methods, albeit with increased computational complexity. Future work will focus on reducing this computational burden and enhancing scalability by incorporating decentralized feedback control laws based on time-varying control barrier functions. Additionally, exploring conflicting temporal logic specifications and other temporal logic languages will enable the application of the framework in dynamically changing environments.

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